

Relational Algebra - Division
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Consider the following tables:

Boats (bid, bname, color)
 Sailors (sid, sname, rating, age)
 Reserves (sid, bid, day)

Division can be represented as \div or $/$
 Ex. $r \div s$ or r / s .

First thing we need to do when we have a division is to consider the **schema**.

| r | |
|---|---|
| x | y |
| | |

| s |
|---|
| y |
| |

| $r \div s$ |
|------------|
| x |
| |

Reserves $\div (\pi_{sid}(\text{Sailors}))$

Second thing we need to consider is **Semantics**. What kind of tuples do we have? Semantics for any x values in r , if x is associated with all y in s , this x value is in $r \div s$.

Ex.

| r | |
|----------------|----------------|
| x | y |
| x ₁ | y ₁ |
| x ₃ | y ₂ |
| x ₂ | y ₃ |

| s |
|----------------|
| y |
| y ₁ |
| y ₂ |
| y ₃ |

In table r , x_i can be associated with one or more than one y . If x_i is not associated with all values of y in s , then it should not be listed in the result table.

Ex.

| r | |
|-----------------|-----------------|
| S _{no} | P _{no} |
| s ₁ | p ₁ |
| s ₁ | p ₂ |
| s ₁ | p ₃ |
| s ₁ | p ₄ |
| s ₂ | p ₁ |
| s ₂ | p ₂ |
| s ₃ | p ₂ |
| s ₄ | p ₂ |
| s ₄ | p ₄ |

| s ₁ |
|-----------------|
| P _{no} |
| p ₂ |
| |

| $r \div s_1$ |
|-----------------|
| S _{no} |
| s ₁ |
| s ₂ |
| s ₃ |
| s ₄ |

If s is associated with all P_{no} in s then it can be in the result table

| |
|----------|
| S_2 |
| P_{no} |
| p_2 |
| p_4 |
| |

| |
|--------------|
| $r \div S_2$ |
| S_{no} |
| s_1 |
| s_4 |
| |

| |
|----------|
| S_3 |
| P_{no} |
| p_1 |
| p_2 |
| p_3 |
| p_4 |
| |

| |
|--------------|
| $r \div S_3$ |
| S_{no} |
| s_1 |
| |
| |

- 1) First step is to define the schema
- 2) Find possible values

Simulate integer division:

$C = A/B \rightarrow$ the largest integer such that $C * B \leq A$

In Relational Algebra

$C = A/B \rightarrow$ the largest relation (# of tuples) such that $C \times B \subseteq A$

Division is not part of the complete set of relational algebraic operators. Then how do we translate a division into an expression that only contains the basic operators?

$r \div s \equiv ?$

Step 1) Find those x values in r that do not satisfy the “all association” conditions

Step 2) $\pi_x(R)$ – result of step 1 (this is a set difference)

It follows the following steps to reach the solution:

Step 1) $\pi_x((\pi_x(r) \times s) - r) \rightarrow$ this step helps find items in $\pi_x(r)$ that don't belong to $r \div s$

Step 2) $\pi_x(r)$ – result from step 1

Ex. Find sailors who reserve all boats

$$\pi_{(sid),(bid)}R \div \pi_{(bid)}(B)$$

More exercises of relational algebra:

Q1. Find sailors with rating > 5 and are younger than 35.

$$\sigma_{(rating > 5) \wedge age < 35} (S)$$

Q2. Find the names of (Q1)

$$\pi_{(sname)} (Q1)$$

Q3. Find names of sailors who reserved boat #103

$\pi_{\text{sname}} (\sigma_{\text{bid}=103} (S \bowtie R))$

Q4. Find sailors who reserved a red boat or a green boat.

$\pi_{\text{sid}} (\sigma_{((\text{color}=\text{green}) \vee (\text{color}=\text{red}))} (R \bowtie B))$

You can also use union

$\pi_{\text{sid}} (\sigma_{\text{color}=\text{green}} (R \bowtie B)) \cup \pi_{\text{sid}} (\sigma_{\text{color}=\text{red}} (R \bowtie B))$

Q5. Find sailors who reserved a green boat and a red boat.

Same as Q4 but instead of union, use intersect

$\pi_{\text{sid}} (\sigma_{\text{color}=\text{green}} (R \bowtie B)) \cap \pi_{\text{sid}} (\sigma_{\text{color}=\text{red}} (R \bowtie B))$

However, we cannot do \wedge instead of \vee for Q5 because in a table an entry cannot be green and red at the same time, so using \wedge will not return anything.